

## Werk

**Titel:** Remarks about the Eilenberg-Zilber type decomposition in cosimplicial sets

**Autor:** Ruiz Salguero, C.; Ruiz S., Roberto

**Jahr:** 1978

**PURL:** [https://resolver.sub.uni-goettingen.de/purl?320387429\\_0012|log8](https://resolver.sub.uni-goettingen.de/purl?320387429_0012|log8)

## Kontakt/Contact

[Digizeitschriften e.V.](#)  
SUB Göttingen  
Platz der Göttinger Sieben 1  
37073 Göttingen

✉ [info@digizeitschriften.de](mailto:info@digizeitschriften.de)

*Revista Colombiana de Matemáticas*  
Vol. XII (1978), págs 61-82

REMARKS ABOUT THE EILENBERG-ZILBER  
TYPE DECOMPOSITION IN COSIMPLICIAL SETS

by

C. RUIZ SALGUERO\* and Roberto RUIZ S.

Contents

- §0 Introduction
- §1 Sections and retractions in the category  $\Delta$ .
- §2 Adjoints of morphisms in the category  $\Delta$ .
- §3 Conditions for unicity of the Eilenberg-Zilber type decomposition in co-simplicial sets.
- §4 Stability of interior points for co-degeneracies.

---

\*Author partially supported by the Universidad Pedagógica Nacional, Bogotá.

§0 Introduction. In [1] the authors have studied the conditions over a model  $Y : \Delta \rightarrow \mathcal{A}$  (or more generally  $Y : \delta \rightarrow \mathcal{A}$ ) that guarantee that the functors  $R_Y : \Delta^\circ \mathcal{S} \rightarrow \mathcal{A}$  (the natural extension of  $Y$  which commutes with inductive limits) commutes with finite products. In order to study this situation in the case  $\mathcal{A} = \Delta^\circ \mathcal{S}$  we need to analyse the set theoretical models  $Y : \Delta \rightarrow \mathcal{S}$  and, in particular, we need to have a theorem corresponding in co-simplicial sets to that which in simplicial sets guarantees the Eilenberg-Zilber decomposition lemma.

To the notion of non-degenerate point in simplicial sets corresponds that of interior points in co-simplicial sets. The Eilenberg-Zilber decomposition lemma guarantees that for each simplicial set  $X$ , and each  $y \in X_n$  there exists one and only one pair  $(\sigma, x)$  where  $\sigma$  is an epimorphism of  $\Delta$  and  $x$  is a non degenerate point of  $X$ , such that  $X(\sigma)(x) = y$ . However, for a point  $y \in Y^n$  ( $Y$  a co-simplicial set) the statement corresponding by duality, namely: "there exists one and only one pair  $(\partial, x)$ , with  $\partial$  a monomorphism of  $\Delta$ , and  $x$  an interior point of  $Y$ , such that  $Y(\partial)(x) = y$ ", is not always true.

We have found that this lack of duality has something to do with the following fact: in a sim-

simplicial set  $X$  every point  $x \in X_0$  belongs to a simplicial point of  $X$  (that is to say, a simplicial subset with only one point in each dimension). This is not so for the co-simplicial case ; there are co-simplicial sets which do not even admit a co-simplicial point. One of the objective of this paper is to show that in order that in a co-simplicial set  $Y$  the unicity of the Eilenber-Zilber decomposition be valid, it is necessary and sufficient that  $Y$  does not admit co-simplicial points. To accomplish this, we are forced to establish the dual of the well known theorem which states that if two epimorphisms of  $\Delta$  have the same sections, then they are equal. This is the point on which the unicity of the decomposition of Eilenberg-Zilber is based for simplicial sets. And it is also to this point that the big difference between simplicial and co-simplicial sets arises, if one uses "mono" instead of "epi" and "retraction" instead of "section" the statement immediately above is not valid in  $\Delta$ . The dual version we have proved is the following "retractions criterion" : if two monomorphisms  $\partial, \partial' : [n] \rightarrow [m]$  of  $\Delta$  have the same retractions and are different then  $n = 0$ .

The relation between the non existence of co-simplicial points in  $Y$  and the retractions criterion is summarized by the equivalence of the two next statement. (i)  $Y$  does not have co-simpli

cial points. (ii) If for two monomorphisms  $\partial, \partial'$  of  $\Delta$ , and for some  $x$ ,  $Y(\partial)(x) = Y(\partial')(x)$ , and  $\text{Ret}(\partial) + \text{Ret}(\partial')$  then necessarily  $\partial = \partial'$ , where  $\text{Ret}(\partial)$  is the set of retractions of  $\partial$ .

We give in this paper another property on a model  $Y$  (which happens to be trivial in the standard cases), necessary to study Milnor's relation, and which permits a characterization of the functor  $R_Y : \Delta^{\circ} \mathcal{S} \rightarrow \mathcal{S}$  (cf. [1]). This property has to do with the stability of interior points under co-degeneracies, we are concerned with whether or not in a co-simplicial set  $Y$  one has for each interior point  $y$  of  $Y$  and each epimorphism  $\sigma$  of  $\Delta$  that  $Y(\sigma)(y)$  is itself an interior point. The answer is negative. But, as we shall see the stability and non existence of co-simplicial points are independent properties. In [1] we will complement these two properties in a model  $Y$  in order to make  $R_Y$  commute with finite products.

§1 Sections and Retractions in the Category  $\Delta$ . Recall that if  $f$  and  $s$  are morphisms of  $\Delta$  such that  $f \circ s = \text{identity}$ , then  $f$  is a retraction of  $s$  and  $s$  is a section of  $f$ . We will denote  $\text{Sec}(f)$  (resp.  $\text{Ret}(s)$ ) the set of sections of  $f$  (resp. retractions of  $s$ ). We also recall two facts.

**1.1 Proposition.** (i) Every monomorphism of  $\Delta$  admits a retraction. (ii) Every epimorphism of  $\Delta$  admits a section.

**1.2 Proposition.** (Section Criterion) If  $f$  and  $f'$  are epimorphism of  $\Delta$  and  $\text{Sec}(f) = \text{Sec}(f')$  then  $f = f'$ .

This last statement is a consequence of the following: given an epimorphism  $f : [n] \rightarrow [m]$  and a point  $x \in [n]$ , then there exists a section  $s$  of  $f$  such that  $x \in \text{Im}(s)$ . Later on, using the concept of adjoint function of an arrow  $\Delta$ , we will give another proof of 1.2.

As we anticipated in the introduction the dual of 1.2 does not hold. In fact, the monomorphisms  $\partial^0, \partial^1 : [0] \rightarrow [1]$  admit a unique retraction  $\sigma^0 : [1] \rightarrow [0]$  without being equal. More generally, any two (mono) morphisms  $[0] \rightarrow [n]$  admits as unique retraction the map  $[n] \rightarrow [0]$ . However, these are the only pathological cases in  $\Delta$ . More precisely :

**1.3 Proposition.** (Retraction Criterion) Let  $\partial, \partial' : [n] \rightarrow [m]$  be two monomorphisms for which  $\text{Ret}(\partial) = \text{Ret}(\partial')$ . If  $\partial \neq \partial'$ , then necessarily  $n = 0$ .

**Proof.** 1. We first show that if  $n \neq 0$ , then  $\partial(n) = \partial'(n)$ . Suppose that  $\partial(n) > \partial'(n)$ . Since

$n \neq 0$ , then  $n - 1 \in [n]$ . We define a function  $\sigma : [m] \rightarrow [n]$  in the following way: for  $x \geq \partial(n)$  let  $\sigma(x) = n$ . On the points of  $[\partial(n) - 1]$  we only require  $\sigma$  to be any retraction of  $\partial \uparrow : [n-1] \rightarrow [\partial(n)-1]$  (which exists by 1.1). In particular, it follows that  $\sigma(\partial(n)-1) = n - 1$ . Such a  $\sigma$  can not be a retraction of  $\partial'$ , because  $\partial(n)-1 \geq \partial'(n)$  and so  $\sigma(\partial(n)-1) \geq \sigma \partial'(n)$ . It follows that  $\sigma \partial'(n) \leq n-1$  and thus  $\sigma \partial'(n) \neq n$ .

2. Dually, it can be proved that if  $n \neq 0$ , and the monomorphism  $\partial, \partial' : [n] \rightarrow [m]$  admit the same retractions, then  $\partial(0) = \partial'(0)$ .

3. Suppose that the monomorphisms  $\partial, \partial' : [n] \rightarrow [m]$  admit the same retractions and  $n \neq 0$ . We know that  $\partial'(n) = \partial(n)$ . The restrictions  $\partial \uparrow, \partial' \uparrow : [n-1] \rightarrow [m]$  also admit the same retractions. If  $n-1 = 0$  then by (2.) above:  $\partial \uparrow(n-1) = \partial' \uparrow(n-1)$  and  $\partial = \partial'$ . If  $n-1 \neq 0$  then by (1.) :  $\partial \uparrow(n-1) = \partial' \uparrow(n-1)$ . By recurrence one completes the proof.

**§2 Adjoints of morphisms in the category  $\Delta$ .** Let  $f : [n] \rightarrow [m]$  be a morphism of  $\Delta$ . Since it is an increasing function it is also a functor between the categories associated with the orders of  $[n]$  and  $[m]$ . Consequently, it makes sense to ask if it admits a right (resp. left) adjoint. If so, the adjoint is an increasing function  $g : [m] \rightarrow [n]$

such that for each  $x \in [n]$ , and each  $y \in [m]$  we have :  $f(x) \leq y \iff x \leq g(y)$ . The last condition is equivalent to the following two : (a) for each  $x \in [n]$ ,  $x \leq gf(x)$  ; (b) for each  $y \in [m]$ ,  $fg(y) \leq y$ . These two conditions represent the morphisms of adjointness. If  $f$  admits a right adjoint  $g$ , then  $f$  commutes with  $\sup$  and  $g$  commutes with  $\inf$ . In our case the last property is trivially satisfied because  $[n]$  and  $[m]$  are finite totally ordered sets, thus the condition becomes the increasingness of the functions. Another necessary condition for the existence of a right (resp. left) adjoint of  $f$  is that  $f(0)=0$  (resp.  $f(n) = m$ ). In fact, applying (b) for  $y = 0$  we have  $gf(0) \leq 0$ , thus  $f^{-1}(0) \neq \emptyset$  and  $f(0)=0$ .

**2.1 Proposition.** In order for  $f : [n] \rightarrow [m]$  to admit a right (resp. left) adjoint it is necessary and sufficient that  $f(0) = 0$  (resp.  $f(n) = m$ ). That is to say  $0 \in \text{Im}(f)$  (resp.  $m \in \text{Im}(f)$ ).

**Proof:** It only remains to show that the condition is sufficient. For each  $y \in [m]$  let  $A(y) = \{x \in [n] \mid f(x) \leq y\}$ .  $A(y)$  is non empty, since  $0 \in A(y)$ . Let  $g(y) = \text{Max } A(y)$ . It follows that  $g : [m] \rightarrow [n]$  is in fact a right adjoint of  $f$ . Dually, if  $f(n) = m$  one defines the left adjoint  $h$  by  $h(y) = \text{Min } B(y)$  where  $B(y) = \{x \in [n] \mid f(x) \geq y\}$ .



Notice that the condition  $f(0) = 0$  is equivalent to the one in the MacLane decomposition of  $f : f = \sigma^{i_s} \dots \sigma^{i_1} \sigma^{j_t} \dots \sigma^{j_1}$ ,  $i_1 > 0$ . Dually  $f(n) = m$  is equivalent to  $m > i_s$ .

If  $f : [n] \rightarrow [m]$  is an epimorphism, then it admits a right adjoint, say  $g$ , and a left adjoint, say  $h$ . Both of them are sections of  $f$ , for they are characterized by

$$g(y) = \text{Max } f^{-1}(y) \quad , \quad h(y) = \text{Min } f^{-1}(y) \quad .$$

For example,  $fg(y) = f \text{Max } f^{-1}(y) = \text{Max } f f^{-1}(y) = \text{Max } \{y\} = y$ .

If we are working with general increasing functions between ordered sets, it is also true that if  $f : X \rightarrow Y$  is an epimorphism and it admits a right adjoint  $g$ , then it is given by  $g(y) = \text{Sup } f^{-1}(y)$  and  $g$  is again a section of  $f$ .

Next we use the order of  $\Delta([n], [m])$  to characterize adjointness of epi and monomorphisms of  $\Delta$ . We define  $f \prec g$  if  $f(x) \leq g(x)$  for each  $x \in [n]$ . Evidently, if  $A$  is a non empty subset of  $\Delta([n], [m])$  then the sup and the inf of  $A$  exist in  $\Delta([n], [m])$ . Moreover, if  $f : [n] \rightarrow [m]$  is an epimorphism then the set  $\text{Sec}(f) \subset \Delta([m], [n])$  admits a maximum and  $f$  is a monomorphism, and  $\text{Ret}(f)$  admits a minimum.

Indeed, let  $g = \text{Sup}(\text{Sec}(f))$  thus for each  $x \in [m]$   
 $g(x) = \text{Sup } v(x) = \text{Max } v(x) \quad (v \in \text{Sec}(f))$ . Then  
 $fg(x) = f(\text{Max } v(x)) = \text{Max } f v(x) = \text{Max } \{x\} = x$ .

**2.2 Proposition.** (a) If  $f : [n] \rightarrow [m]$  is an epimorphism then the right adjoint of  $f$  is  $\text{Max}(\text{Sec}(f))$ .

(b) If  $\partial : [n] \rightarrow [m]$  is a monomorphism admitting left adjoint, say  $f$ , then  $f$  is a retraction of  $\partial$  and  $f = \text{Min}(\text{Ret}(\partial))$ .

Proof. (a) Let  $g$  be the right adjoint of  $f$  and  $u = \text{Max}(\text{Sec}(f))$ . Since  $g$  is a section of  $f$ ,  $g \leq u$ . Furthermore, by adjointness,  $x \leq gf(x)$ , thus  $x \leq uf(x)$ . Since  $fu(y) = y$ , for each  $y$ ,  $u$  satisfies properties (a) and (b) of adjointness of  $f$ . Since in  $[n]$  and  $[m]$  the isomorphisms are equalities,  $u = g$ .

(b) For each  $x \in [m]$ ,  $f(x) = \text{Inf}\{y \mid \partial(y) \geq x\}$ . Then  $f\partial(y) = \text{Inf}\{y' \mid \partial(y') \geq \partial(y)\}$ . Since  $\partial$  is a monomorphism this inf is precisely  $y$ . That proves the first statement of part (b). The second one is proven by a similar procedure to that in part (a).

**2.3 Alternative proofs of the retraction and section criteria.** For the retraction criterion : Let

$\partial, \partial' : [n] \rightarrow [m]$  be monomorphisms satisfying  $\text{Ret}(\partial) = \text{Ret}(\partial')$ . We have already seen that if  $n \neq 0$ , then  $\partial(n) = \partial'(n)$ . Let  $\delta, \delta' : [n] \rightarrow \rightarrow [\partial(n)]$  denote the functions obtained from  $\partial$  and  $\partial'$  by codomain restriction. Then  $\delta$  and  $\delta'$  admit left adjoints and  $\text{Ret}(\delta) = \text{Ret}(\delta')$ . Since  $\text{Min Ret}(\delta) = \text{Min Ret}(\delta')$ , then by 2.2 the left adjoint of  $\delta$  coincides with that of  $\delta'$ . Thus  $\delta = \delta'$  and also  $\partial = \partial'$ .

For the section criterion, contrary to the retraction criterion, the proof is direct, for if two epimorphisms  $\sigma, \sigma'$  have the same set of sections then both admit right adjoint and  $\text{ad}(\sigma) = \text{Max Sec}(\sigma) = \text{Max Sec}(\sigma') = \text{ad}(\sigma')$ . So  $\sigma = \sigma'$ .

### §3 Conditions for the unicity of the Eilenberg-Zilber type decomposition in co-simplicial sets.

**3.1 Definition.** Let  $Y : \Delta \rightarrow \mathcal{S}$  be a co-simplicial set and let  $y \in Y^n = Y([n])$ . We say that  $y$  is interior, or  $y$  is an interior point of  $Y$ , if the following condition holds "if there exist  $p \geq 0$ , a monomorphism  $\partial : [p] \rightarrow [n]$ , and  $y' \in Y^p$ , such that  $Y(\partial)(y') = y$ , then  $p = n$  and  $\partial = 1_{[n]}$ ". In other words  $y$  is an interior point of  $Y$  if either  $y \in Y^o$ , or  $y \in Y^n$  with

$n > 0$  and  $y$  does not belong to the image of the co-faces  $Y(\partial^i)$   $i = 0, \dots, n$ .

It is clear that for a point  $y \in Y^n$  there are two possibilities: either there exist a monomorphism  $\partial : [m] \rightarrow [n]$  which is not an isomorphism such that  $y \in \text{Im}(Y(\partial))$ , or every monomorphism  $\partial$  for which  $y \in \text{Im}(Y(\partial))$  is an isomorphism hence the identity. In the latter case,  $y$  is an interior point.

Now, if  $y$  is not an interior point, it can be written in the form  $y = Y(\partial)(y')$  with  $\partial$  a monomorphism, and so  $\dim y' < \dim y = n$ . If  $y'$  is not an interior point then  $y' = Y(\partial')(y'')$ ; therefore,  $y = Y(\partial\partial')(y'')$ . This process can always be continued until an interior point  $z$  and a monomorphism  $\delta$  are found such that  $y = Y(\delta)(z)$ .

**3.2 Lemma-Definition.** For each  $y \in Y^n$  ( $Y$  a co-simplicial set) there always exist a monomorphism  $\delta$  in  $\Delta$  and an interior point  $z$  of  $Y$  such that  $y = Y(\delta)(z)$ . In such a case, the pair  $\langle \delta, z \rangle$  is called an Eilenberg-Zilber type decomposition of  $y$  (E-Z decomposition).

We emphasize that, contrary to what happens in simplicial sets, in general the E-Z co-simplicial decomposition is not unique. In fact, if  $Y^n$  has only one point for each  $n$ , then the point

$x_1 \in Y^1$  is written in to different ways  $x_1 = Y(\partial^0)(x_0) = Y(\partial^1)(x_0)$ . Moreover, the only co-simplicial sets  $Y$  in which there are points with more than one E-Z decomposition are (as we shall see) those in which there exists a point  $x_0$  in  $Y^0$  such that  $Y(\partial^0)(x_0) = Y(\partial^1)(x_0)$ ,  $(\partial^0, \partial^1 : [0] \rightarrow [1])$ . Actually, the E-Z decompositions of a point have common characteristics which reveal the properties needed by a model  $Y$  in order to have the "unique E-Z decomposition" property. We think of these properties as a kind of partial uniqueness and devote our next proposition to them.

**3.3 Proposition.** Let  $\partial, \partial'$  be monomorphism of  $\Delta$  and  $y, y'$  interior points of  $Y$ . If  $Y(\partial)(y) = Y(\partial')(y')$ , then (i)  $y = y'$  and (ii)  $\text{Ret}(\partial) = \text{Ret}(\partial')$ .

**Proof.** Let  $\sigma : [n] \rightarrow [m]$  (resp  $\sigma' : [n] \rightarrow [m']$ ) be a retraction of  $\partial : [m] \rightarrow [n]$  (resp  $\partial' : [m'] \rightarrow [n]$ ), whose existence was already proven. Mapping the identity  $Y(\partial)(y) = Y(\partial')(y')$  by  $Y(\sigma)$ , we get that  $y = Y(\sigma\partial')(y')$  since  $\sigma\partial = 1_{[m]}$ . Using the MacLane decomposition  $\sigma\partial' = \delta\circ\mu$  where  $\delta$  is a monomorphism and  $\mu$  is an epimorphism, we get  $y = Y(\delta)(Y(\mu)(y'))$ . Since  $y$  is interior and  $\delta$  is a monomorphism,  $\delta$  is an identity and consequently  $\mu = \sigma\partial' : [m'] \rightarrow [m]$  is an

epimorphism. Thus  $m' \geq m$ . With the same kind of procedure one shows that  $m \geq m'$ . Hence  $m = m'$ . Since the only epimorphism  $[m] \rightarrow [m]$  is the identity one gets  $\sigma\partial' = 1_{[m]}$  and  $\sigma'\partial = 1_{[m]}$ . Thus we have proven that (i)  $y = Y(\sigma\delta')(y') = Y(\text{id})(y') = y'$  and (ii)  $\text{Ret}(\partial) = \text{Ret}(\partial')$ .

Remark: The proof just presented corresponds in the cosimplicial case to the one presented by Gabriel and Zisman in [2] for simplicial sets, on which ours was inspired.

3.4 Corollary. If  $Y(\partial)(y) = Y(\partial')(y')$ , where  $\partial : [m] \rightarrow [n]$  and  $\partial' : [m'] \rightarrow [n]$  are monomorphisms of  $\Delta$ , and  $y, y'$  are interior points of  $Y$ , then : (i)  $m = m'$ , (ii)  $y = y'$ , (iii) if  $m \neq 0$  then  $\partial = \partial'$ .

The proof of this corollary is an immediate consequence of the retraction criterion (1.3). Notice also that when  $m = 0$  we cannot conclude that  $\partial = \partial'$ , but (iii) can be put in a more suggestive way: (iii') if  $\partial \neq \partial'$  then  $\partial, \partial' : [0] \rightarrow [n]$ .

3.5 Definition. (1) A co-simplicial set  $Y$  is said to be of the Eilenberg-Zilber type (E-Z type) if every  $y \in Y$  has a unique E-Z decomposition.

(2) A co-simplicial set  $Y$  admits a co-sim-

a co-simplicial point if there exists a co-simplicial subset of  $Y$  with exactly one point in each dimension.

**3.6 Lemma.** In order for  $y \in Y^0$  to be an element of a co-simplicial point of  $Y$  it is necessary and sufficient that  $Y(\partial^0)(y) = Y(\partial^1)(y)$  ( $\partial^0, \partial^1: [0] \rightarrow [1]$ ).

Proof. That the condition is necessary is clear. The sufficiency follows by induction on  $n$ . If  $\partial, \partial' : [0] \rightarrow [n]$  are two arrows of  $\Delta$ , then  $Y(\partial)(y) = Y(\partial')(y)$  (which would imply that  $y$  belongs to a co-simplicial point of  $Y$ ). In fact, for  $n = 1$  it is the hypothesis. Assume it holds for  $k < n$  and let  $\partial, \partial' : [0] \rightarrow [n]$ . For  $\partial$  (and  $\partial'$ ) there are two possibilities  $\partial(0) = n$ , or  $\partial(0) \neq n$ . In other words  $\partial = \partial^{n-1} \circ \delta$  or  $\partial = \partial^n \circ \delta$  for some  $\delta : [0] \rightarrow [n]$  (also  $\partial' = \partial^{n-1} \circ \delta$  or  $\partial' = \partial^n \circ \delta$  where  $\delta' : [0] \rightarrow [n-1]$ ).

From the four possibilities there are two which follow directly by induction hypothesis. As the other two are treated similarly, we present only one case, say  $Y(\partial^n \delta)(y) = Y(\partial^{n-1} \delta')(y)$ . Let  $\mu = \partial^{n-1} \circ \dots \circ \partial^1 : [0] \rightarrow [n-1]$ . By the induction hypothesis  $Y(\mu)(y) = Y(\delta)(y) = Y(\delta')(y)$ . Then  $Y(\partial^n) Y(\delta)(y) = Y(\partial^n) Y(\mu)(y) = Y(\partial^n \partial^{n-1} \dots \partial^1)$ . Similarly  $Y(\partial^{n-1}) Y(\delta')(y) = Y(\partial^{n-1}) Y(\mu)(y) =$

$= Y(\partial^{n-1} \partial^{n-1} \dots \partial^1)(y) = Y(\partial^n \partial^{n-1} \dots \partial^1)(y)$  because  $\partial^{n-1} \partial^{n-1} = \partial^n \partial^{n-1}$ . This ends the proof.

**3.7 Lemma.** In order that  $Y$  admit a co-simplicial point it is necessary and sufficient that there exists two different arrows  $\partial, \partial' : [0] \rightarrow [n]$  and  $y \in Y^0$  such that  $Y(\partial)(y) = Y(\partial')(y)$ .

Proof. The condition is evidently necessary. Conversely we will prove by induction on  $k$  the proposition  $P(k) : "$  if there exist different arrows  $\partial, \partial' : [0] \rightarrow [k]$  and  $y \in Y^0$  such that  $Y(\partial)(y) = Y(\partial')(y)$ , then the co-simplicial set  $Y$  admits a co-simplicial point".  $P(1)$  is the previous lemma. Suppose  $P(k)$  for  $k < n$ . Let's prove  $P(n)$ . Using the same technique as in 3.6,  $\partial = \partial^n \circ \delta$  or  $\partial = \partial^{n-1} \circ \delta$  for some  $\delta : [0] \rightarrow [n-1]$ . Similarly,  $\partial' = \partial^n \circ \delta'$  or  $\partial' = \partial^{n-1} \circ \delta'$ ,  $\delta' : [0] \rightarrow [n-1]$ . In either case we apply  $Y(\sigma^{n-1})$  to the identity  $Y(\partial)(y) = Y(\partial')(y)$ , from which we get the existence of  $\delta, \delta' : [0] \rightarrow [n-1]$  such that  $Y(\delta)(y) = Y(\delta')(y)$ . If  $\delta \neq \delta'$ , we apply the induction hypothesis to find a co-simplicial point, but if  $\delta = \delta'$  we cannot use the induction hypothesis. In that case, one has  $\partial = \partial^n \circ \delta$ ,  $\partial' = \partial^{n-1} \circ \delta$  (resp.  $\partial = \partial^{n-1} \circ \delta$ ,  $\partial' = \partial^n \circ \delta$ ) since  $\partial \neq \partial'$ . The MacLane decomposition of  $\partial'$  must be  $\partial' = \partial^{n-1} \partial^{n-1} \dots \partial^0$ . We apply  $Y(\sigma^{n-2})$ ,



coface which exists because  $n \geq 2$ , to the equality  $Y(\partial)(y) = Y(\partial')(y)$  obtaining  $Y(\partial^{n-1} \partial^{n-3} \dots \partial^0)(y) = Y(\partial^{n-2} \partial^{n-3} \dots \partial^0)(y)$ . But  $\partial^{n-1} \partial^{n-3} \dots \partial^0 \neq \partial^{n-2} \partial^{n-3} \dots \partial^0$  (MacLane decomposition), and now we may apply the induction hypothesis.

**3.8 Theorem.** For a co-simplicial set  $Y$  the following statements are equivalent: (1)  $Y$  does not admit co-simplicial points. (2)  $Y$  is an E-Z type co-simplicial set. (3) For any pair of morphisms  $\partial, \partial' : [p] \rightarrow [n]$  such that  $\text{Ret}(\partial) = \text{Ret}(\partial')$ , if there exist  $x \in Y^p$  for which  $Y(\partial)(x) = Y(\partial')(x)$  then  $\partial = \partial'$ .

Proof. (2)  $\Rightarrow$  (1) is evident. (1)  $\Rightarrow$  (3) since otherwise there would exist  $\partial, \partial' : [p] \rightarrow [n]$  with  $\text{Ret}(\partial) = \text{Ret}(\partial')$ ,  $\partial \neq \partial'$  and  $x \in Y^p$  such that  $Y(\partial)(x) = Y(\partial')(x)$ . By the retraction criterion (1.3),  $p = 0$ . By the previous lemma,  $Y$  admits cosimplicial points. Finally, (3)  $\Rightarrow$  (2): suppose that  $z$  has two E-Z decompositions, say  $z = Y(\partial)(x) = Y(\partial')(x')$ . Then by (3.3)  $x = x'$ ,  $\text{Ret}(\partial) = \text{Ret}(\partial')$  and, by hypothesis,  $\partial = \partial'$ . consequently  $\langle x, \partial \rangle = \langle x', \partial' \rangle$ .

**§4 Stability of interior points under co-degeneracies.** In a cosimplicial set, if  $y \in Y^n$  is an

interior and  $\sigma : [n] \rightarrow [m]$  is an epimorphism then  $Y(\sigma)(y)$  is not necessarily interior. In other words, it may happen that  $Y(\sigma)(x) = Y(\partial)(x')$  with  $\sigma$  an epimorphism,  $\partial$  a monomorphism and  $x, x'$  interior points, but the arrows being non trivial. It is our purpose to exhibit co-simplicial sets with this feature and to observe that the property of being of E-Z type is not enough to make it disappear.

Take, for example, a simplicial set  $X$  which in dimension 2 has two different non degenerate points  $a$  and  $b$  such that  $d_0(a) = d_1(a) = d_2(a) = d_0(b) = d_1(b) = d_2(b)$ . That is the case with  $K(G, 2)$  or more generally with any simplicial group  $K$  for which  $\Pi_2(K) \neq 0$ . Let  $C$  be a "sufficiently large" set. Let  $v : X_0 \rightarrow C$  be a function, and  $w = v \circ d_0 = v \circ X(\partial^0) : X_1 \rightarrow C$ . We define  $u : X_2 \rightarrow C$  as follows:  $u(s_0(x)) = w(x)$  for any  $x \in X_1$ . For  $a$  and  $b$  above, we take  $u(a)$  and  $u(b)$  to be two different points of  $C$ . For the other points of  $X_2$  it does not matter how  $u$  is defined. We denote by  $Y$  the cosimplicial set with  $Y^n = \mathcal{S}(X_n, C)$ , and co-faces induced by faces of  $X$  by composition. The point  $u \in Y^2$  cannot be factored through  $d_0, d_1, d_2 : X_2 \rightarrow X_1$  and therefore it is interior. On the other hand,  $Y(\sigma^0)(u) = s_0 \circ u = w = v \circ d_0 = Y(\partial^0)(v)$  and therefore it is not an interior point.

We now give some examples of co-simplicial sets with stable interior points

4.1 Definition. A co-simplicial set is said to satisfy M0.2 (cf [1]) if for every  $n \geq 0$ , every interior point  $x \in Y^n$  and every epimorphism  $\sigma : [n] \rightarrow [p]$ ,  $Y(\sigma)(x)$  is also interior point.

Examples:

4.2 Let  $p \geq 0$  and  $Y(\ ) = \Delta([p], -) : \Delta \rightarrow \mathcal{S}$ . A point  $x : [p] \rightarrow [n]$  is interior when it is an epimorphism of  $\Delta$ . It is evident that if  $\sigma : [n] \rightarrow [m]$  is an isomorphism then  $\sigma \circ x = Y(\sigma)(x)$  is also an interior point. This model does not have co-simplicial points. Notice that in terms of the E-Z property this means that in  $\Delta$  any arrow  $\alpha : [p] \rightarrow [n]$  is decomposable in the form  $\partial \circ \sigma$  where  $\partial$  is a mono and  $\sigma$  an epimorphism, and this decomposition is unique. That is to say, the E-Z type decomposition of these models ( $p \geq 0$ ) is equivalent to the unique Mac-Lane decomposition in  $\Delta$ .

4.3 The co-simplicial set  $\Delta(\ ) : \Delta \rightarrow \mathcal{S}$  defined by  $\Delta(n) = \{(t_0, \dots, t_n) \mid 0 \leq t_i \leq 1, \sum t_i = 1\}$ . If  $\alpha : [n] \rightarrow [m]$  then  $\Delta(\alpha)(x) = (T_0, \dots, T_m)$ , where  $x = (t_0, \dots, t_n)$  and  $T_i = \sum t_j$ , the sum running over the set  $\{j \mid \alpha(j) = i\}$ . When this last set is empty,  $T_i = 0$ . In

this co-simplicial set a point  $x = (t_0, \dots, t_n)$  is interior if none of the  $t_i$ 's is zero. Evidently  $\Delta(\alpha)(x)$  is also interior if and only if  $\alpha$  is an epimorphism. Notice also that this model does not have cosimplicial points.

4.4 The co-simplicial set  $\mathcal{P}_0(\ ) : \Delta \rightarrow \mathcal{S}$  which associates to each  $[n]$  the set of non empty parts of  $[n] = \{0, 1, \dots, n\}$ , and to each  $\alpha : [m] \rightarrow [n]$  the map  $\mathcal{P}_0(\alpha) =$  direct image by  $\alpha$ . In this case a point  $A \in \mathcal{P}_0([n])$  is interior if and only if  $A = [n]$ . This characteristic is certainly preserved by epimorphisms. Since we have eliminated the empty set from the set of parts, this model does not have co-simplicial points and consequently is an E-Z co-simplicial set. The unicity of the E-Z decomposition becomes simply the fact that a totally finite ordered set can be enumerated in only one way respecting its order and beginning at zero. In this example as in the others,  $Y^0$  is a point.

4.5 More generally, for each integer  $p \geq 0$  let  $\Delta'[\ ]_p : \Delta \rightarrow \mathcal{S}$  be the co-simplicial set given for each  $n$  by  $\Delta'[n]_p = \{(A_0, \dots, A_p) \mid \emptyset \neq A_0 \subseteq A_1 \subseteq \dots \subseteq A_p \subseteq [n]\}$ , and for each  $\alpha : [n] \rightarrow [m]$  by  $\Delta'[\alpha]_p(A_0, \dots, A_p) = (\alpha(A_0), \dots, \alpha(A_p))$ . In this case  $(A_0, \dots, A_p)$  is interior in dimension

$n$  if and only if  $\Delta_p = [n]$ . This property is again preserved by epimorphisms. Moreover, if  $\alpha$  preserves one interior point then  $\alpha$  must be an epimorphism. Example 4.5 is simply the  $p$ -th dimension of Kan's first sub-division over  $\Delta[n]$ . The model  $\Delta'[\ ]_p$  do not have co-simplicial points and the E-Z decomposition of  $x = (A_0, \dots, A_p)$  with  $\emptyset \neq A_0 \subseteq \dots \subseteq A_p \subseteq [n]$  can be given in the following simple way. Let  $q = \text{card}(A_p) - 1$ ; there exist one and only one monotone map  $\alpha : [q] \rightarrow [n]$  such that  $\alpha([q]) = A_p$ . We define  $B_i = \alpha^{-1}(A_i)$ , thus  $\Delta'[\alpha]_p (B_0, \dots, B_p) = (A_0, \dots, A_p)$ . The properties M0.2 and E-Z of these co-simplicial sets are used in [1] in order to prove that Kan's first sub-division does not commute with finite products.

**4.6 Remark.** In our examples the property M0.2 and the non existence of co-simplicial points are present together. That is not true in general. In fact, if in example 4.5 we drop the condition " $A_i \neq \emptyset$ " and denote the co-simplicial set by  $Y_p$ , then the element of  $Y_p^\circ$  of the form  $(A_0, \dots, A_p)$  with  $A_j = \emptyset$  for every  $j$  is the only one which generates a co-simplicial point. However, a point  $y = (A_0, \dots, A_p)$  is interior if  $\dim(y) = 0$  or if  $\dim(y) = n > 0$  and  $A_p = [n]$ . Thus,  $Y_p$  has property M0.2.

4.7 Remark. We now face the inverse of situation 4.6. That is to say, we will provide an example of a E-Z co-simplicial set  $Y$  which fails to have M0.2 . We will take the example at the beginning of the present section (§4) which, as we know , fails to have both M0.2 and E-Z properties. We then exhibit a procedure which allows us to eliminate the co-simplicial points. We then make sure that this procedure does not eliminate the M0.2 failure.

If a co-simplicial set  $A$  has co-simplicial points then one can get from it a co-simplicial set without co-simplicial points by eliminating all the points which by some co-face co-degeneracy fall into a co-simplicial point. A characterization of the eliminated points can be given as follows : let  $x \in Y^P$  , then "there exist  $\varepsilon: [p] \rightarrow [m]$  such that  $Y(\varepsilon)(x)$  belongs to a co-simplicial point if and only if  $Y(\eta)(x)$  belongs to a co-simplicial point, where  $\eta: [p] \rightarrow [0]$  ". We recall that  $Y(\eta)(x)$  belongs to a co-simplicial point if and only if  $Y(\partial^0 \eta)(x) = Y(\partial^1 \eta)(x)$ . If in our example, at the beginning of the section, we do the surgery just described, it remains to see that if the point  $v$  is not a co-simplicial point then it is not eliminated. In fact, if it were eliminated then  $Y(\partial^0 \eta)(u) = Y(\partial^1 \eta)(u)$  for  $\eta = \sigma^0 \sigma^0: [2] \rightarrow [0]$ . Since by construction  $Y(\partial^0)(v) = Y(\sigma^0)(u)$  , one gets  $Y(\partial^0)(v) = Y(\partial^1)(v)$ .

BIBLIOGRAPHY

- [1] C. Ruiz-Salguero and R. Ruiz, Conditions over a "realization" functor in order for it to commute with finite products (to appear).
- [2] P. Gabriel and M. Zisman, Calculus of fractions and homotopy theory, *Ergebnisse der Mathematik und Ihrer Grenzgebiete*, Band 35, Springer-Verlag 1967.

\*\*\*

*Departamento de Matemáticas*  
*Universidad Nacional de Colombia*  
*Bogotá, D.E., COLOMBIA.*

*Departamento de Matemáticas*  
*Universidad del Valle.*  
*Cali, Valle, COLOMBIA.*

(Recibido en octubre de 1978)