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LOCAL DETERMINACY OF SYMMETRIC PSEUDOPROCESSES

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A number of various physical systems can be described by means of relations in the cartesian product $P \times T$, where P is the set of all possible states of the system concerned and T is a set of time instants. This approach to the study of the behaviour of systems was used in [1] to [5], where a wide class of such relations is investigated in detail. The present paper is a direct continuation of [5] so that as far as the notation and the terminology is concerned, the reader is referred to [5]. To make the text of the paper as self-contained as possible, the basic notions and notation from [5] will be recalled in the first point of the next section.

1. SYMMETRIC PSEUDOPROCESSES

1.1. Notation. In what follows, P denotes an arbitrary set, R the set of all reals, $R^* = R \cup \{+\infty, -\infty\}$ the extended real line with the ordering extended from R to R^* in the natural way, T a subset of R .

If X, Y are sets, then any subset of the cartesian product $X \times Y$ (in this order) is called a relation between X and Y . If $X = Y$, then a relation $r \subset X \times X$ is called a relation in X . The relation inverse to a relation r is denoted by r^{-1} . The identity relation in X is denoted by 1_X . If $r \subset X \times Y$, $s \subset Y \times Z$, then the composition of the relations r and s (in this order) is denoted by $r \circ s$. If a pair $(x, y) \in X \times Y$ belongs to a relation $r \subset X \times Y$, then we write either $(x, y) \in r$ or xry . Given $r \subset X \times Y$, we set

$$(1.1.1) \quad D_r = \{y \in Y \mid xry \text{ for some } x \in X\},$$

$$(1.1.2) \quad I_r = \{x \in X \mid xrx\} \text{ if } X = Y,$$

$$(1.1.3) \quad ry = \{x \in X \mid (x, y) \in D_r\},$$

$$(1.1.4) \quad rA = \{x \in X \mid (x, y) \in D_r \text{ for some } y \in A\},$$

$$(1.1.5) \quad xr = \{y \in Y \mid (x, y) \in D_r\},$$

$$(1.1.6) \quad Br = \{y \in Y \mid (x, y) \in D_r \text{ for some } x \in B\},$$

$$(1.1.7) \quad r|_A = r \cap (X \times A),$$

for each $y \in Y, A \subset Y, x \in X, B \subset X$.

In the present paper we shall be concerned mainly with relations p, q in $P \times T$, i.e. with subsets of $(P \times T) \times (P \times T)$. Each such relation $q \subset (P \times T) \times (P \times T)$ can be uniquely described by the two-parametric system of relations ${}_v q_u$ in P with $u, v \in T$ as follows:

$$(1.1.8) \quad (y, v) q(x, u) \text{ iff } y {}_v q_u x, \quad x, y \in P, \quad u, v \in T.$$

A relation p in $P \times T$ such that

$$(I) \quad {}_u p_u \subset 1_P \text{ for each } u \in T$$

and

$$(R) \quad {}_v p_u \neq \emptyset \text{ implies } u \leq v \text{ for all } u, v \in T$$

is called a *right pseudoprocess in P over T* . The set of all right pseudoprocesses in P over T is denoted by $\text{Ps}(P, T)$. A right pseudoprocess $p \in \text{Ps}(P, T)$ is said to be a *compositive right pseudoprocess*, a *transitive right pseudoprocess* or a *right process in P over T* iff the condition

$$(RC) \quad {}_v p_u \subset {}_v p_t \circ {}_t p_u \text{ for all } u \leq t \leq v \text{ in } T,$$

$$(RT) \quad {}_v p_u \supset {}_v p_t \circ {}_t p_u \text{ for all } u \leq t \leq v \text{ in } T$$

or

$$(RP) \quad {}_v p_u = {}_v p_t \circ {}_t p_u \text{ for all } u \leq t \leq v \text{ in } T,$$

is satisfied, respectively. The set of all compositive right pseudoprocesses, transitive right pseudoprocesses and right processes in P over T will be denoted by $\text{Psc}(P, T)$, $\text{Pst}(P, T)$ and $P(P, T)$. A more detailed explanation of the theory of right pseudoprocesses may be found in [5].

1.2. Definition. Let P be an arbitrary set, $T \subset R, q$ a relation in $P \times T$. The relation q is called a *symmetric pseudoprocess in P over T* iff it satisfies the conditions

$$(I) \quad {}_u q_u \subset 1_P \text{ for all } u \in T,$$

$$(S) \quad {}_v q_u = ({}_u q_v)^{-1} \text{ for all } u, v \in T.$$

The set of all symmetric pseudoprocesses in P over T will be denoted by $\text{Ss}(P, T)$.

1.3. Remark. The property (S) in 1.2 may be reformulated as

$$(1.3.1) \quad y_v q_u x \text{ iff } x_u q_v y \text{ for all } x, y \in P, u, v \in T.$$

Hence, this property is equivalent with

$$(1.3.2) \quad q = q^{-1}.$$

Thus

$$(1.3.3) \quad {}_v q_u = {}_u q_v \text{ for all } u, v \in T.$$

For a symmetric pseudoprocess q , the sets D_q and I_q from (1.1.1) and (1.1.2) may be characterized as follows:

$$D_q = \{(x, u) \in P \times T \mid {}_v q_u x \neq \emptyset \text{ for some } v \in T\}$$

and

$$I_q = \{(x, u) \in D_q \mid x_u q_u x\}.$$

1.4. Construction. Let $p \in \text{Ps}(P, T)$. It is not difficult to verify that the relation $p \cup p^{-1}$ in $P \times T$ fulfils the conditions of Definition 1.2 so that it is a symmetric pseudoprocess in P over T . The symmetric pseudoprocess q in P over T defined by

$$(1.4.1) \quad q = p \cup p^{-1}$$

is said to be *induced by the right pseudoprocess* p in P over T .

Let us show that, given a symmetric pseudoprocess $q \in \text{Ss}(P, T)$, there exists a right pseudoprocess $p \in \text{Ps}(P, T)$ such that (1.4.1) holds.

1.5. Definition. Let $q \in \text{Ss}(P, T)$, $q^+ \in \text{Ps}(P, T)$. The right pseudoprocess q^+ is said to be *positively induced by the symmetric pseudoprocess* q iff it satisfies the condition

$$(1.5.1) \quad {}_v q^+_u = {}_v q_u \text{ for all } u \leq v \text{ in } T.$$

1.6. Remark. Since q^+ is a right pseudoprocess, it holds

$$(1.6.1) \quad {}_v q^+_u = \emptyset \text{ for all } u > v \text{ in } T$$

so that we obtain from 1.5 and (1.3.1) that

$$(1.6.2) \quad D_{q^+} \subset D_q.$$

The inclusion in (1.6.2) cannot be in general replaced by the equality. However, if $q \in \text{Ss}(P, T)$ is such that for each $(x, u) \in D_q$ there exists $t \in T$ fulfilling the conditions $t \geq u$ and ${}_t q_u x \neq \emptyset$, then equality

$$(1.6.3) \quad D_{q^+} = D_q$$

holds.

1.7. Lemma. Let $q, \tilde{q} \in \text{Ss}(P, T)$. Then the following assertions hold:

- (i) $q = q^+ \cup (q^+)^{-1}$.
- (ii) $q \subset \tilde{q}$ iff $q^+ \subset \tilde{q}^+$.

1.8. Definition. Let $p \in \text{Ps}(P, T)$, $q \in \text{Ss}(P, T)$, $p' \in \text{Ps}(P, -T)$, $q' \in \text{Ss}(P, -T)$, where $-T = \{t \in R \mid -t \in T\}$.

The right pseudoprocess p' is said to be *orientation-change produced from the right pseudoprocess p* iff

$$(1.8.1) \quad {}_v p'_u = (-{}_u p_{-v})^{-1} \text{ for all } u \leq v \text{ in } -T.$$

The symmetric pseudoprocess q' is said to be *orientation-change produced from the symmetric pseudoprocess q* iff

$$(1.8.2) \quad {}_v q'_u = (-{}_u q_{-v})^{-1} \text{ for all } u, v \text{ in } -T.$$

1.9. Definition. Let $q \in \text{Ss}(P, T)$, let $q^+ \in \text{Ps}(P, T)$ be positively induced by q , and let $q^- \in \text{Ps}(P, -T)$.

The right pseudoprocess q^- is said to be *negatively induced* by the symmetric pseudoprocess q iff q^- is orientation-change produced from q^+ .

1.10. Remark. A right pseudoprocess q^- from the preceding definition can be described directly by the relations ${}_v q_u$ as follows.

The equality

$$(1.10.1) \quad {}_v q^-_u = (-{}_u q^+_{-v})^{-1} = (-{}_u q_{-v})^{-1} = -{}_v q_{-u}$$

holds for all $u \leq v$ in $-T$, hence

$$(1.10.2) \quad {}_v q_u = -{}_v q^-_{-u} \text{ for all } v \leq u \text{ in } T.$$

Clearly

$$(1.10.3) \quad D_{q^-} \subset D_{q'} = \{(x, u) \in P \times (-T) \mid (x, -u) \in D_q\},$$

where the symmetric pseudoprocess q' is orientation-change produced from the symmetric pseudoprocess q . If for each $(x, u) \in D_q$ there exists $t \leq u$ in T such that ${}_t q_u x \neq \emptyset$, then the inclusion in (1.10.3) can be replaced by the equality.

1.11. Lemma. Let $q, \tilde{q} \in \text{Ss}(P, T)$ and let $q' \in \text{Ss}(P, -T)$ be orientation-change produced from q . Then the following assertions hold:

- (i) $q' = q^- \cup (q^-)^{-1}$.
- (ii) $q \subset \tilde{q}$ iff $q^- \subset \tilde{q}^-$.

Proof. The assertions follow from (1.8.2), 1.2 (ii), (1.10.1), (1.10.2) and from Lemmas 1.2 and 1.3 in [5].

1.12. Corollary. Let $q, \sim q \in \text{Ss}(P, T)$. Then the following three inclusions are equivalent:

- (i) $q \subset \sim q$;
- (ii) $q^+ \subset \sim q^+$;
- (iii) $q^- \subset \sim q^-$.

1.13. Definition. Let $q \in \text{Ss}(P, T)$, $s \subset P \times T$. The relation s is called a *solution* of the symmetric pseudoprocess q iff the following three conditions are satisfied:

- (i) the domain D_s of s is an interval in T ;
- (ii) s is a map of D_s into P ;
- (iii) $s(v) \circ q_u s(u)$ holds for all $u, v \in D_s$.

The set of all solutions of q will be denoted by S_q .

1.14. Theorem. Let $q, \sim q \in \text{Ss}(P, T)$, let $q' \in \text{Ss}(P, -T)$ be orientation-change produced from q and let $q^+ \in \text{Ps}(P, T)$ and $q^- \in \text{Ps}(P, -T)$ be the right pseudoprocesses positively and negatively induced by q , respectively. Let $s : T \rightarrow P$ and $s' : -T \rightarrow P$ be maps such that D_s is an interval in T ; $D_{s'} = \{t \mid -t \in D_s\}$ and $s'(-t) = s(t)$ for all $t \in D_s$. Then the following assertions hold:

- (i) $s \in S_q$ iff $s \times s \subset q$.
- (ii) $s \in S_q$ iff $s' \in S_{q'}$.
- (iii) $S_q = S_{q^+}$.
- (iv) $S_{q'} = S_{q^-}$.
- (v) $S_{q \cap \sim q} = S_q \cap S_{\sim q}$.

1.16. Definition. Let $q \in \text{Ss}(P, T)$. The maps

$$(1.16.1) \quad e^+ : D_q \rightarrow R^{\#}, \quad e^- : D_q \rightarrow R^{\#}$$

defined by

$$(1.16.2) \quad e^+(x, u) = \sup \{t \in T \mid {}_t q_u x \neq \emptyset\},$$

$$(1.16.3) \quad e^-(x, u) = \inf \{t \in T \mid {}_t q_u x \neq \emptyset\}$$

are called the *positive* and the *negative extent of existence* of q , respectively.

1.17. Remark. Let us recall that if q^+ or q^- is a right pseudoprocess positively or negatively induced by a symmetric pseudoprocess q , then the extents of existence e or e' of these pseudoprocesses are defined, according to Definition 2.3 in [5], by

$$(1.17.1) \quad e(x, u) = \sup \{t \in T \mid {}_t q^+_u x \neq \emptyset\}, \quad (x, u) \in D_{q^+}$$

or

$$(1.17.2) \quad e'(x, u) = \sup \{t \in T \mid {}_t q^-_u x \neq \emptyset\}, \quad (x, u) \in D_{q^-}.$$

If $(x, u) \in D_{q^+}$, then $(x, u) \in D_q$ so that $e^+(x, u)$ as well as $e(x, u)$ are defined and it is evident that the equality

$$(1.17.3) \quad e^+(x, u) = e(x, u), \quad (x, u) \in D_{q^+}$$

holds. Analogously, if $(x, -u) \in D_{q^-}$, then $(x, u) \in D_q$ so that both $e'(x, -u)$ and $e^-(x, u)$ are defined and the equality

$$(1.17.4) \quad e^-(x, u) = -e'(x, -u)$$

takes place. From (1.16.2) and (1.16.3) we obtain immediately the inequality

$$(1.17.5) \quad e^-(x, u) \leq e^+(x, u) \quad \text{for all } (x, u) \in D_q.$$

1.18. Definition. Let $q \in \text{Ss}(P, T)$.

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) local existence at a point* $(x, u) \in D_q$ iff $e^+(x, u) > u$ ($e^-(x, u) < u$, $e^-(x, u) < u < e^+(x, u)$).

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) local existence* iff it has positive (negative, bilateral) local existence at each point $(x, u) \in D_q$.

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) global existence at a point* $(x, u) \in D_q$ iff $e^+(x, u) = \sup T$ ($e^-(x, u) = \inf T$, $e^+(x, u) = \sup T$ and $e^-(x, u) = \inf T$).

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) global existence* iff it has positive (negative, bilateral) global existence at each point $(x, u) \in D_q$.

If $q^+ \in \text{Ps}(P, T)$ is a right pseudoprocess positively induced by a symmetric pseudoprocess q , then q^+ is said to have some of the properties described above iff q has the property.

1.19. Definition. Let $q \in \text{Ss}(P, T)$. A point $(x, u) \in I_q$ is called a *start point* or an *end point* of the symmetric pseudoprocess q iff ${}_t q_u x = \emptyset$ holds for all $t \in T$ such that $t < u$ or $u < t$, respectively.

1.20. Remark. Let a (right or symmetric) pseudoprocess r in P over T be given. In accordance with the notation introduced in [5], item 5.1, the symbol \mathcal{S}_r or \mathcal{E}_r will stand for the set of all start or end points of the pseudoprocess r , respectively.

Let $q \in \text{Ss}(P, T)$, $(x, u) \in D_q$. Then $(x, u) \in \mathcal{S}_q$ or $(x, u) \in \mathcal{E}_q$ iff $e^-(x, u) = u$ or $e^+(x, u) = u$, respectively. If $(x, u) \in \mathcal{S}_q$ or $(x, u) \in \mathcal{E}_q$ and $s \in S_q$ is such that $s(u) = x$, then $u = \min D_s$ or $u = \max D_s$, respectively. The converse of this assertion is not valid.

1.21. Lemma. Let $q \in \text{Ss}(P, T)$. Let $q^+ \in \text{Ps}(P, T)$ or $q^- \in \text{Ps}(P, -T)$ be positively or negatively induced by q , respectively. Let $(x, u) \in I_q$. Then the following assertions hold.

- (i) $\mathcal{S}_q = \mathcal{S}_{q^+}$, $\mathcal{E}_q = \mathcal{E}_{q^+}$.
(ii) $(x, u) \in \mathcal{S}_q$ iff $(x, -u) \in \mathcal{E}_{q^-}$, $(x, u) \in \mathcal{E}_q$ iff $(x, -u) \in \mathcal{S}_{q^-}$.

Proof follows from 1.2 (ii), (1.5.1) and (1.10.2).

1.22. Definition. Let $q \in \text{Ss}(P, T)$. The maps

$$(1.22.1) \quad d^+ : D_q \rightarrow R^* \quad \text{and} \quad d^- : D_q \rightarrow R^*$$

defined by

$$(1.22.2) \quad d^+(x, u) = \sup \{w \in R \mid \text{card}({}_t q_u x) \leq 1 \text{ for all } t \in T \cap \langle u, w \rangle\},$$

$$(1.22.3) \quad d^-(x, u) = \inf \{w \in R \mid \text{card}({}_t q_u x) \leq 1 \text{ for all } t \in T \cap \langle w, u \rangle\}$$

are called the *positive* and the *negative extent of unicity* of the symmetric pseudoprocess q .

1.23. Remark. Notice that $+\infty$ may belong to the range of the function d^+ and that d^+ can assume this value also in the case of a bounded T . Similarly for d^- and $-\infty$. In general, it holds

$$(1.23.1) \quad -\infty \leq d^-(x, u) \leq u \leq d^+(x, u) \leq +\infty.$$

Let q^+ and q^- be the right pseudoprocesses positively and negatively induced by a symmetric pseudoprocess q , respectively. Then, according to Definition 2.5 in [5] the extents of unicity d and d' of q^+ and q^- are defined by

$$(1.23.2) \quad d(x, u) = \sup \{w \in R \mid \text{card}({}_t q^+_u x) \leq 1 \text{ for all } t \in T \cap \langle u, w \rangle\}$$

for all $(x, u) \in D_{q^+}$,

and

$$(1.23.3) \quad d'(x, u) = \sup \{w \in R \mid \text{card}({}_t q^-_u x) \leq 1 \text{ for all } t \in (-T) \cap \langle u, w \rangle\}$$

for all $(x, u) \in D_{q^-}$,

respectively. Then for each $(x, u) \in D_{q^+} \subset D_q$ both $d(x, u)$ and $d^+(x, u)$ are defined and

$$(1.23.4) \quad d^+(x, u) = d(x, u) \quad \text{for all } (x, u) \in D_{q^+}.$$

Analogously, for each $(x, -u) \in D_{q^-}$ it holds $(x, u) \in D_q$ so that $d'(x, -u)$ as well as $d^-(x, u)$ are defined and the equality

$$(1.23.5) \quad d^-(x, u) = -d'(x, -u) \quad \text{for all } (x, -u) \in D_{q^-}$$

takes place.

1.24. Definition. Let $q \in \text{Ss}(P, T)$.

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) local unicity at a point* $(x, u) \in D_q$ iff $d^+(x, u) > u$ ($d^-(x, u) < u$, $d^-(x, u) < u < d^+(x, u)$).

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) local unicity* iff it has positive (negative, bilateral) local unicity at each point $(x, u) \in D_q$.

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) global unicity at a point* $(x, u) \in D_q$ iff $d^+(x, u) = +\infty$ ($d^-(x, u) = -\infty$, $d^+(x, u) = -d^-(x, u) = +\infty$).

The symmetric pseudoprocess q is said to have *positive (negative, bilateral) global unicity* iff it has positive (negative, bilateral) global unicity at each point $(x, u) \in D_q$.

If $q^+ \in \text{Ps}(P, T)$ is a right pseudoprocess positively induced by a symmetric pseudoprocess q , then q^+ is said to have some of the properties described above iff q has the property.

1.25. Lemma. Let $q \in \text{Ss}(P, T)$. Then the following assertions hold:

- (i) If $(x, u) \in D_q$ is a start or an end point of q , then $d^-(x, u) = -\infty$ or $d^+(x, u) = +\infty$, respectively.
- (ii) If $(x, u) \in D_q$ and $u < d^+(x, u) < +\infty$ or $-\infty < d^-(x, u) < u$, then $d^+(x, u) < e^+(x, u)$ or $e^-(x, u) < d^-(x, u)$, respectively.

1.26. Lemma. Let $q, \tilde{q} \in \text{Ss}(P, T)$, $q \subset \tilde{q}$. Let e^+, e^-, d^+, d^- and $\tilde{e}^+, \tilde{e}^-, \tilde{d}^+, \tilde{d}^-$ be the corresponding extents of existence and unicity of q and \tilde{q} , respectively. Then the following assertions hold:

- (i) If $(x, u) \in D_q$, then

$$(1.26.1) \quad e^+(x, u) \leq \tilde{e}^+(x, u), \quad \tilde{e}^-(x, u) \leq e^-(x, u),$$

$$(1.26.2) \quad \tilde{d}^+(x, u) \leq d^+(x, u), \quad d^-(x, u) \leq \tilde{d}^-(x, u).$$

- (ii) If $(x, u) \in D_{\tilde{q}}$ is a start point or an end point of \tilde{q} , then it is a start point or an end point of q , respectively.
- (iii) If q has positive, negative, bilateral local or global existence at a point $(x, u) \in D_q$, then \tilde{q} has the same property.
- (iv) If \tilde{q} has positive, negative, bilateral local or global unicity at a point $(x, u) \in D_{\tilde{q}}$, then q has the same property.

1.27. Definition. Let $q \in \text{Ss}(P, T)$. The symmetric pseudoprocess q is said to be *solution complete* iff for each pair $((y, v), (x, u)) \in q$ there exists $s \in S_q$ such that $s(u) = x$, $s(v) = y$.

1.28. Theorem. Let $q \in Ss(P, T)$. Then the symmetric pseudoprocess q is solution complete iff the right pseudoprocess q^+ positively induced by q is solution complete.

Proof follows directly from 1.7 (i), 1.14 (iii) and 1.27.

2. SYMMETRIC PROCESSES

2.1. Definition. Let $q \in Ss(P, T)$ and let $q^+ \in Ps(P, T)$ be positively induced by q .

The symmetric pseudoprocess q is said to be *compositive* or *transitive* iff the right pseudoprocess q^+ is compositive or transitive, respectively.

The symmetric pseudoprocess q is called a *symmetric process in P over T* iff it is compositive and transitive.

The sets of all compositive symmetric pseudoprocesses, of all transitive symmetric pseudoprocesses and of all symmetric processes in P over T will be denoted by $Ssc(P, T)$, $Sst(P, T)$ and $S(P, T)$, respectively.

2.2. Lemma. Let $q \in Ss(P, T)$, let $q^+ \in Ps(P, T)$ be positively induced by q and let $q^- \in Ps(P, -T)$ be negatively induced by q . Then the following assertions are equivalent:

- (i) $q \in Ssc(P, T)$.
- (ii) $q^+ \in Psc(P, T)$.
- (iii) $q^- \in Psc(P, -T)$.
- (iv) ${}_v q_u \subset {}_v q_t \circ {}_t q_u$ for all $u \leq t \leq v$ in T .
- (v) ${}_v q_u \subset {}_v q_t \circ {}_t q_u$ for all $v \leq t \leq u$ in T .
- (vi) ${}_v q_u \subset {}_v q_t \circ {}_t q_u$ for all $u, v, t \in T$, t between u, v .

Proof follows from 2.1, (1.10.2), (RC) in 1.1, (1.5.1) and (1.10.4).

2.3. Lemma. Let $q \in Ss(P, T)$, let $q^+ \in Ps(P, T)$ be positively induced by q and let $q^- \in Ps(P, -T)$ be negatively induced by q . Then the following assertions are equivalent:

- (i) $q \in Sst(P, T)$.
- (ii) $q^+ \in Pst(P, T)$.
- (iii) $q^- \in Pst(P, -T)$.
- (iv) ${}_v q_t \circ {}_t q_u \subset {}_v q_u$ for all $u \leq t \leq v$ in T .
- (v) ${}_v q_t \circ {}_t q_u \subset {}_v q_u$ for all $v \leq t \leq u$ in T .
- (vi) ${}_v q_t \circ {}_t q_u \subset {}_v q_u$ for all $u, v, t \in T$, t between u, v .

2.4. Lemma. Let $q \in \text{Ss}(P, T)$. Then the following assertions are equivalent:

- (i) $q \in \text{S}(P, T)$.
- (ii) ${}_v q_u = {}_v q_t \circ {}_t q_u$ for all $u, v, t \in T$, t between u, v .

2.5. Lemma. If $q \in \text{Ssc}(P, T)$, then

$$D_{q^+} = D_q = I_q,$$

$$D_{q^-} = \{(x, u) \in P \times (-T) \mid (x, -u) \in D_q\}.$$

2.6. Lemma. Let $q \in \text{Ss}(P, T)$ and let I be an arbitrary set. Then the following assertions hold:

- (i) If $s \in S_q$, J an interval in T , then $s|_J \in S_q$.
- (ii) If $s_i \in S_q$ for $i \in I$ are such that $D_{\cap s_i}$ is an interval in T , then $\cap s_i \in S_q$.
- (iii) If q is transitive and if $s_i \in S_q$ with $i \in I$ are such that $D_{s_i} \cap D_{s_j} \neq \emptyset$ and $s_i \cup s_j$ is a map for all $i, j \in I$, then $s_i \in S_q$.

2.7. Lemma. Let $q \in \text{Sst}(P, T)$, $(x, u) \in D_q$. Then the following assertions hold:

- (i) If $v, w \in T$, v between u, w , $z_w q_v y$, $y_v q_u x$, then also $z_w q_u x$.
- (ii) Let $(y, v) \in D_q$ be such that $y_v q_u x$. If $v = e^-(x, u)$ or $v = e^+(x, u)$, then (y, v) is a start point or an end point of q , respectively.

2.8. Lemma. Let $q \in \text{Ssc}(P, T)$, $(x, u) \in D_q$, $u, v, w \in T \cap \langle d^-(x, u), d^+(x, u) \rangle$, v between u, w . If $y_v q_u x$, $z_w q_u x$, then also $z_w q_v y$.

2.9. Theorem. Let $q \in \text{Ssc}(P, T)$ have global unicity and let $s : T \rightarrow P$. Then $s \in S_q$ iff the following two conditions are satisfied:

- (i) D_s is an interval in T ;
- (ii) there exists $u \in D_s$ such that $s(v) {}_v q_u s(u)$ holds for all $v \in D_s$.

Proof follows easily from Definition 1.13 and Lemma 2.8.

2.10. Theorem. Let $q \in \text{Ss}(P, T)$. If q is solution complete, then it is compositive.

Proof. According to Theorem 1.28 the right pseudoprocess q^+ is solution complete so that it is compositive (see Theorem 3.8 in [5]). Now apply Definition 2.1.

2.11. Definition. Let $q \in \text{Ss}(P, T)$, $(x, u) \in D_q$, $s \subset P \times T$. The relation s is called a *characteristic solution of q through the point (x, u)* iff it satisfies the following two conditions:

- (i) $D_s = \{v \in T \mid \text{card}({}_t q_u x) = 1 \text{ for all } t \in T, t \text{ between } u, v\}$;
- (ii) $s(v) {}_v q_u x$ holds for all $v \in D_s$.

2.12. Lemma. Let $q \in \text{Ssc}(P, T)$ have bilateral local unicity at a point $(x, u) \in D_q$ and let s be the characteristic solution of q through (x, u) . Then $s \in S_q$ with

$$\begin{aligned} \sup D_s &= \min \{e^+(x, u), d^+(x, u)\} = \begin{cases} e^+(x, u) & \text{if } d^+(x, u) = +\infty, \\ d^+(x, u) & \text{if } d^+(x, u) < +\infty, \end{cases} \\ \inf D_s &= \max \{e^-(x, u), d^-(x, u)\} = \begin{cases} e^-(x, u) & \text{if } d^-(x, u) = -\infty, \\ d^-(x, u) & \text{if } d^-(x, u) > -\infty. \end{cases} \end{aligned}$$

2.13. Remark. In Definition 3.10 in [5], each right pseudoprocess p in P over T is associated with the maximal compositive right pseudoprocess \hat{p} in P over T contained in p , the so called lower modification of p . The construction of the lower modification of a right pseudoprocess described in item 3.12 in [5] is applicable to symmetric pseudoprocesses as well. However, this is not necessary, because, as will be shown, the maximal compositive symmetric pseudoprocess contained in a given symmetric pseudoprocess q , which will be called again the lower modification of q , can be constructed directly from the lower modifications of the pertinent right pseudoprocesses q^+ and q^- .

2.14. Theorem. Let $q \in \text{Ss}(P, T)$, let $q^+ \in \text{Ps}(P, T)$ be positively induced by q , let $q^- \in \text{Ps}(P, T)$ be negatively induced by q , let \hat{q}^+ be the lower modification of q^+ and let \hat{q}^- be orientation-change produced from \hat{q}^+ . Then \hat{q}^- is the lower modification of q^- .

Proof. According to Definition 1.8, $\hat{q}^- \in \text{Ps}(P, -T)$. First we shall prove that \hat{q}^- is compositive.

The right pseudoprocess \hat{q}^+ being the lower modification of q^+ is compositive, i.e.

$$(2.14.1) \quad {}_v \hat{q}^+_u \supset {}_v \hat{q}^+_t \circ {}_t \hat{q}^+_u \text{ for all } u \leq t \leq v \text{ in } T.$$

Hence

$$(2.14.2) \quad ({}_v \hat{q}^+_u)^{-1} \supset ({}_v \hat{q}^+_t \circ {}_t \hat{q}^+_u)^{-1} \text{ for all } u \leq t \leq v \text{ in } T.$$

According to (1.10.1) it is

$$(2.14.3) \quad ({}_v \hat{q}^+_u)^{-1} = {}_{-u} \hat{q}^-_{-v}$$

so that

$$(2.14.4) \quad ({}_v \hat{q}^+_t \circ {}_t \hat{q}^+_u)^{-1} = ({}_t \hat{q}^+_u)^{-1} \circ ({}_v \hat{q}^+_t)^{-1} = {}_{-u} \hat{q}^-_{-t} \circ {}_{-t} \hat{q}^-_{-v}$$

holds for all $u \leq t \leq v$ in T . Substituting from (2.14.3) and (2.14.4) into (2.14.2) we obtain

$${}_{-u} \hat{q}^-_{-v} \subset {}_{-u} \hat{q}^-_{-t} \circ {}_{-t} \hat{q}^-_{-v} \text{ for all } -v \leq -t \leq -u \text{ in } -T.$$

Thus \hat{q}^- is compositive.

Now we shall prove that \hat{q}^- is the maximal compositive right pseudoprocess in P over $-T$ contained in q^- .

Let $p \in \text{Psc}(P, -T)$ be such that $\hat{q}^- \subset p \subset q^-$. Then

$$(2.14.5) \quad {}_v \hat{q}^-_u \subset {}_v p_u \subset {}_v q^-_u \quad \text{for all } u \leq v \text{ in } -T.$$

Denote by p' the right pseudoprocess in P over T which is orientation-change produced from p . From (2.14.5) and (1.10.1) one obtains

$$(-_u \hat{q}^+_{-v})^{-1} \subset -_u p'_{-v} \subset -_u q^+_{-v} \quad \text{for all } -v \leq -u \text{ in } T,$$

which can be written equivalently as

$$\hat{q}^+ \subset p' \subset q^+.$$

Since \hat{q}^+ and p' are compositive, it is necessarily $\hat{q}^+ = p'$, hence $\hat{q}^- = p$.

We have proved that \hat{q}^- is the lower modification of q^- .

2.15. Theorem. Let $q \in \text{Ss}(P, T)$, let $q^+ \in \text{Ps}(P, T)$ be positively induced by q and let \hat{q}^+ be the lower modification of q^+ . Then the symmetric pseudoprocess

$$(2.15.1) \quad \hat{q} = \hat{q}^+ \cup (\hat{q}^+)^{-1}$$

is the maximal compositive symmetric pseudoprocess in P over T contained in q .

Proof. Since \hat{q}^+ is compositive, \hat{q} is compositive as well.

Let $\tilde{q} \in \text{Ssc}(P, T)$ be such that $\hat{q} \subset \tilde{q} \subset q$. Then

$$(\hat{q}^+ \cup (\hat{q}^+)^{-1}) \subset (\tilde{q}^+ \cup (\tilde{q}^+)^{-1}) \subset (q^+ \cup (q^+)^{-1}).$$

Hence, according to 1.12, one obtains $\hat{q}^+ \subset \tilde{q}^+ \subset q^+$. Since \hat{q}^+ is the lower modification of q^+ and \tilde{q}^+ is compositive, it is necessarily $\hat{q}^+ = \tilde{q}^+$, hence $\hat{q} = \tilde{q}$ follows by virtue of Lemma 1.12.

2.16. Definition. Let $q \in \text{Ss}(P, T)$. The symmetric compositive pseudoprocess \hat{q} in P over T defined by

$$(2.16.1) \quad \hat{q} = \hat{q}^+ \cup (\hat{q}^+)^{-1}$$

is called the *lower modification* of the symmetric pseudoprocess q .

2.17. Remark. One may verify easily that

$$D_{\hat{q}} = D_{\hat{q}^+} \subset D_{q^+} \subset D_q.$$

If, in addition, $D_q = I_q$, then

$$D_{\hat{q}} = D_{\hat{q}^+} = D_{q^+} = D_q.$$

2.18. Theorem. Let $q_i \in \text{Ss}(P, T)$ for $i = 1, 2$ and let q_i^+ be the right pseudoprocess positively induced by q_i . Then $q_1 \cap q_2 \in \text{Ss}(P, T)$ and its lower modification is the compositive symmetric pseudoprocess $q_1 \wedge q_2$ in P over T defined by

$$(2.18.1) \quad q_1 \wedge q_2 = (q_1^+ \wedge q_2^+) \cup (q_1^+ \wedge q_2^+)^{-1},$$

where $q_1^+ \wedge q_2^+$ denotes the lower modification of the right pseudoprocess $q_1^+ \cap q_2^+$.

Proof. According to the assertion (i) of Lemma 1.7 it is

$$q_1 = q_1^+ \cup (q_1^+)^{-1}, \quad q_2 = q_2^+ \cup (q_2^+)^{-1}$$

so that

$$q_1 \cap q_2 = (q_1^+ \cap q_2^+) \cup (q_1^+ \cap q_2^+)^{-1}.$$

Hence and from 1.7 (i) one obtains

$$(2.18.2) \quad (q_1 \cap q_2)^+ = q_1^+ \cap q_2^+.$$

Since the lower modification of $q_1^+ \cap q_2^+$ is $q_1^+ \wedge q_2^+$, the equality (2.18.1) follows now directly from (2.18.2) and (2.16.1).

2.19. Theorem. Let $q \in \text{Ss}(P, T)$, let $\wedge q$ be its lower modification. Then $S_{\wedge q} = S_q$.

Proof. Applying Theorem 6.14 to the equalities

$$q = q^+ \cup (q^+)^{-1}, \quad \wedge q = \wedge q^+ \cup (\wedge q^+)^{-1}$$

we obtain $S_q = S_{q^+}$, $S_{\wedge q} = S_{\wedge q^+}$. Theorem 3.14 in [5] yields $S_{\wedge q^+} = S_{q^+}$. Thus $S_{\wedge q} = S_q$.

2.20. Corollary. Let $q_1, q_2 \in \text{Ss}(P, T)$. Then $S_{q_1 \wedge q_2} = S_{q_1 \cap q_2}$.

3. LOCAL BEHAVIOUR OF SYMMETRIC PSEUDOPROCESSES

3.1. In Section 5 of the paper [5] we have investigated the local behaviour of right pseudoprocesses. Let us recall the basic notions and notation which will be used in the sequel.

Given a (right or symmetric) pseudoprocess r in P over T , the symbol L_r will denote the set

$$(3.1.1) \quad L_r = \{(s, u) \in S_r \times T \mid u \in D_s\}.$$

Let $p \in \text{Ps}(P, T)$, $(x, u) \in D_p$. Then p is said to have right (or left) local existence of solutions at the point (x, u) iff the following conditions are fulfilled:

- (i) $(x, u) \notin \mathcal{E}_p$ (or $(x, u) \notin \mathcal{S}_p$);
- (ii) There exist $\varepsilon > 0$ and $s \in S_p$ such that

$$\langle u, u + \varepsilon \rangle \cap T \subset D_s \quad (\text{or } \langle u - \varepsilon, u \rangle \cap T \subset D_s).$$

A pseudoprocess p is said to have bilateral local existence of solutions at the point (x, u) iff it has right local existence of solutions at the point (x, u) if $(x, u) \in D_p - \mathcal{E}_p$ and left local existence of solutions at the point (x, u) if $(x, u) \in D_p - \mathcal{S}_p$. A pseudoprocess p is said to have right or left or bilateral local existence of solutions iff it has the property at each point $(x, u) \in D_p$.

Let $p, \tilde{p} \in \text{Ps}(P, T)$. Then \tilde{p} is said to determine the local behaviour of p (which is shortly written as $\tilde{p} < p$) iff $\tilde{p} \subset p$ and there exists a map

$$(3.1.2) \quad k : L_p \rightarrow R$$

such that $k(s, u) > u$ for $u < \sup D_s$, $k(s, u) = u$ for $u = \max D_s$ and

$$(3.1.3) \quad s|_{\langle u, k(s, u) \rangle} \in S_{\tilde{p}}.$$

A pseudoprocess \tilde{p} is said to determine the bilateral local behaviour of p (which is shortly written as $\tilde{p} \leq p$) iff $\tilde{p} \subset p$ and there exist maps

$$(3.1.4) \quad k_1, k_2 : L_p \rightarrow R$$

such that

$$k_1(s, u) < u \text{ for } \inf D_s < u, \quad k_1(s, u) = u \text{ for } \min D_s = u,$$

$$k_2(s, u) > u \text{ for } \sup D_s > u, \quad k_2(s, u) = u \text{ for } \max D_s = u$$

and

$$(3.1.5) \quad s|_{\langle k_1(s, u), k_2(s, u) \rangle} \in S_{\tilde{p}}.$$

Now, let $q \in \text{Ss}(P, T)$ and $q^+ \in \text{Ps}(P, T)$ be positively induced by q . According to Theorem 1.14 (iii) it holds $S_q = S_{q^+}$. This enables us to define the corresponding notions related to the local existence of solutions and to the local behaviour of pseudoprocesses for symmetric pseudoprocesses in a natural way as follows.

3.2. Definition. Let $q \in \text{Ss}(P, T)$, let $q^+ \in \text{Ps}(P, T)$ be positively induced by q and let $(x, u) \in D_q$. The symmetric pseudoprocess q is said to have *right* or *left* or *bilateral local existence of solutions at the point* (x, u) iff the right pseudoprocess q^+ has right or left or bilateral local existence of solutions at the point (x, u) , respectively.

The symmetric pseudoprocess q is said to have *right* or *left* or *bilateral local existence of solutions* iff it has the property at each point $(x, u) \in D_q$.

3.3. Definition. Let $q \in \text{Ss}(P, T)$, $p, \tilde{p}, q^+ \in \text{Ps}(P, T)$, $\tilde{p}, q^- \in \text{Ps}(P, -T)$, where q^+ is positively and q^- negatively induced by q . The right pseudoprocess \tilde{p}, \tilde{p} or p is said to *determine the negative local behaviour, the positive local behaviour or the local behaviour* of the symmetric pseudoprocess q (which is shortly written as $\tilde{p} \leq q, \tilde{p} < q$ or $p \leq q$) iff $\tilde{p} < q^-, \tilde{p} < q^+$, or $p \leq q^+$.

Let $q, \tilde{q} \in \text{Ss}(P, T)$. The symmetric pseudoprocess \tilde{q} is said to *determine the local behaviour* of the symmetric pseudoprocess q (shortly written as $\tilde{q} \leq q$) iff $\tilde{q}^+ \leq q$, where \tilde{q}^+ is the right pseudoprocess positively induced by \tilde{q} .

3.4. Remark. If $q, \tilde{q} \in \text{Ss}(P, T)$, $\tilde{q} \leq q$, then $D_{\tilde{q}} \subset D_q$. If, in addition, $D_q = I_q$, then $D_{\tilde{q}} = D_q$. This equality holds in particular if $q \in \text{Ssc}(P, T)$.

Since Definitions 3.2 and 3.3 are immediate generalizations of Definitions 5.2, 5.3, 5.5 and 5.6 from [5], it is natural that many results valid for right pseudoprocesses remain valid as well when formulated for symmetric pseudoprocesses. Some results of this kind are given in what follows.

Similarly as in [5] many assertions concerning the relations $<$ and \leq may be formulated simultaneously. It will be done using the symbol $<$. In these assertions the symbol $<$ has to be replaced either by $<$ or by \leq .

3.5. Lemma. Let $q \in \text{Ss}(P, T)$, $p \in \text{Ps}(P, T)$, $\tilde{p} \in \text{Ps}(P, -T)$. Then the following assertions hold:

- (i) If $p < q$, then $S_p \subset S_q$, $D_p \subset D_q$, $\mathcal{E}_p = \mathcal{E}_q \cap D_p$.
- (ii) If $\tilde{p} \leq q$, then $S_{\tilde{p}} \subset S_{q^-}$.
- (iii) If $p \leq q$, then $\mathcal{E}_p = \mathcal{E}_q \cap D_p$, $\mathcal{L}_p = \mathcal{L}_q \cap D_p$.
- (iv) If $p < q$ and p has right or left local existence of solutions at each point, then $D_p = D_q$, $\mathcal{E}_p = \mathcal{E}_q$; if, in addition, $p \leq q$, then also $\mathcal{L}_p = \mathcal{L}_q$.

3.6. Lemma. Let $q, \tilde{q}, \tilde{\tilde{q}} \in \text{Ss}(P, T)$. If $q \approx \tilde{\tilde{q}} \leq \tilde{q}$, $\tilde{q} \leq q$, then $\tilde{\tilde{q}} \leq q$.

3.7. Lemma. Let $p, \tilde{p} \in \text{Ps}(P, T)$, $q \in \text{Ss}(P, T)$, $\tilde{p} \subset p \subset q$. If $\tilde{p} < q$, then $p < q$.

3.8. Lemma. Let $p \in \text{Ps}(P, T)$, $q, \tilde{q} \in \text{Ss}(P, T)$, $p \subset \tilde{q} \subset q$. If $p < q$, then $p < \tilde{q}$.

3.9. Lemma. Let $p \in \text{Ps}(P, T)$, $q \in \text{Ss}(P, T)$. Then $p < q$ iff $\hat{p} < q$, where \hat{p} is the lower modification of p .

3.10. Lemma. Let $p, \tilde{p} \in \text{Ps}(P, T)$, $q \in \text{Ss}(P, T)$. Then the following assertions are equivalent.

- (i) $p < q$, $\tilde{p} < q$.
- (ii) $p \cap \tilde{p} < q$.
- (iii) $p \wedge \tilde{p} < q$.

3.11. Lemma. Let $q, \tilde{q}, \tilde{\tilde{q}} \in \text{Ss}(P, T)$. If $\tilde{q} \leq q$, then $\tilde{q} \cap q \approx \tilde{\tilde{q}} \leq q \cap \tilde{\tilde{q}}$. Especially, if $\tilde{q} \leq q$, $\tilde{\tilde{q}} \leq q$, then $\tilde{q} \cap \tilde{\tilde{q}} \leq \tilde{q}$, $\tilde{q} \cap \tilde{\tilde{q}} \leq \tilde{\tilde{q}}$.

3.12. Lemma. Let $q \in \text{Ssc}(P, T)$, $\sim q, \approx q \in \text{Ss}(P, T)$, $q \preceq \sim q, q \preceq \approx q$. Then $q^+ \subset \subset \sim q^+ \wedge \approx q^+, q^- \subset \subset \sim q^- \wedge \approx q^-, q \subset \subset \sim q \wedge \approx q$.

3.13. Theorem. Let $q \in \text{Ss}(P, T)$, $p \in \text{Ps}(P, T)$ and let $p' \in \text{Ps}(P, -T)$ be orientation change produced from p . Then the following assertions hold:

(i) $p' \preceq q$ iff $p \subset q$ and there exists a map $h^- : L_q \rightarrow R$ such that

$$h^-(s, u) < u \text{ for } \inf D_s < u, \quad h^-(s, u) = u \text{ for } \min D_s = u$$

and

$$(3.13.1) \quad s(t) \text{ } {}_t p_v s(v) \text{ for all } h^-(s, u) \leq v \leq t \leq u \text{ in } D_s.$$

(ii) $p < q$ iff $p \subset q$ and there exists a map $h^+ : L_q \rightarrow R$ such that

$$h^+(s, u) > u \text{ for } u < \sup D_s, \quad h^+(s, u) = u \text{ for } u = \max D_s$$

and

$$(3.13.2) \quad s(t) \text{ } {}_t p_v s(v) \text{ for all } u \leq v \leq t \leq h^+(s, u) \text{ in } D_s.$$

(iii) $p \preceq q$ iff $p \subset q$ and there exist maps $h^+, h^- : L_q \rightarrow R$ such that

$$h^+(s, u) > u \text{ for } u < \sup D_s, \quad h^+(s, u) = u \text{ for } u = \max D_s,$$

$$h^-(s, u) < u \text{ for } \inf D_s < u, \quad h^-(s, u) = u \text{ for } u = \min D_s$$

and

$$(3.13.3) \quad s(t) \text{ } {}_t p_v s(v) \text{ for all } h^-(s, u) \leq v \leq t \leq h^+(s, u) \text{ in } D_s.$$

Proof. Before proving the assertion (i) let us recall that (1.8.1) and (1.10.1) yield $p' \subset q^-$ iff $p \subset q^+$ and according to 1.14 (ii), $s \in S_q = S_{q^+}$ iff there exists $s' \in S_{q^-}$ such that $s'(t) = s(-t)$ for all $-t \in D_s$.

Suppose $p' \preceq q$, i.e. $p' \subset q^-$ and prove that (3.13.1) is fulfilled. Take $(s, u) \in L_q$ arbitrary. According to the assumption there exists a real $h'(s', -u)$ such that

$$(3.13.4) \quad s'(-v) \text{ } {}_{-v} p'_{-t} s'(-t) \text{ for all } -u \leq -t \leq -v \leq h'(s', -u) \text{ in } D_s.$$

Hence, setting $h^-(s, u) = -h'(s', -u)$ and using (1.8.1), one easily obtains (3.13.1).

Suppose now that the condition (3.13.1) is fulfilled and prove that $p' \subset q^-$. The condition (3.13.1) can be written in the form

$$(3.13.5) \quad s(-t) \text{ } {}_{-t} p_{-v} s(-v) \text{ for all } h^-(s, u) \leq -v \leq -t \leq u \text{ with } v, t \in D_s,$$

i.e.

$$s'(v) \text{ } {}_v p'_t s'(t) \text{ for all } -u \leq t \leq v \leq -h^-(s, u) \text{ with } v, t \in D_s.$$

Setting $h'(s', u) = -h^-(s, -u)$ we conclude that for each $(s', u) \in L_{q^-}$ there exists a real $h'(s', u)$ such that

$$s'(t) {}_t p_v s'(v) \text{ for all } u \leq v \leq t \leq h'(s', u) \text{ in } D_{s'}.$$

Thus $p' \prec q^-$.

The assertions (ii) and (iii) follow immediately from Definition 3.3.

3.14. Theorem. *Let $q, \sim q \in \text{Ss}(P, T)$. Then the following three assertions are equivalent:*

- (i) $\sim q \preceq q$;
- (ii) $\sim q^+ \preceq q^+$;
- (iii) $\sim q^- \preceq q^-$.

Proof follows easily from Theorem 1.14 and Definition 3.3.

3.15. Theorem. *Let $q \in \text{Ss}(P, T)$, $\sim q \in \text{Sst}(P, T)$. If $\sim q^+ \prec q$, $\sim q^- \preceq q$, then $\sim q \preceq q$.*

Proof. According to 1.12 the inclusion $\sim q \subset q$ is equivalent to any one of the inclusions $\sim q^+ \subset q^+$ and $\sim q^- \subset q^-$.

To each $(s, u) \in L_q$ we can assign reals $h^+(s, u)$ and $h^-(s, u)$ as in Theorem 3.13 such that

$$(3.15.1) \quad s(t) {}_t \sim q_v s(v) \text{ for all } u \leq v \leq t \leq h^+(s, u) \text{ in } T$$

and

$$(3.15.2) \quad s(t) {}_t \sim q_v s(v) \text{ for all } h^-(s, u) \leq v \leq t \leq u \text{ in } T.$$

Especially,

$$s(t) {}_t \sim q_u s(u), \quad s(u) {}_u \sim q_v s(v) \text{ for all } h^-(s, u) \leq v \leq u \leq t \leq h^+(s, u) \text{ in } T.$$

Since $\sim q$ is transitive, it holds also

$$(3.15.3) \quad s(t) {}_t \sim q_v s(v) \text{ for all } h^-(s, u) \leq v \leq u \leq t \leq h^+(s, u) \text{ in } T.$$

Finally, $({}_t \sim q_v)^{-1} = {}_v \sim q_t$, so that $s(v) {}_v \sim q_t s(t)$ holds iff $s(t) {}_t \sim q_v s(v)$. This together with (3.15.1), (3.15.2) and (3.15.3) yields

$$s(t) {}_t \sim q_v s(v) \text{ for all } v, t \in D_s \cap \langle h^-(s, u), h^+(s, u) \rangle,$$

which was to be proved.

3.16. Definition. Let $q, \tilde{q} \in \text{Ss}(P, T)$. The symmetric pseudoprocesses q and \tilde{q} are said to be *negatively locally equivalent* or *positively locally equivalent* or *locally equivalent* (which is shortly written as $q \preceq \succeq \tilde{q}$ or $q \prec \succ \tilde{q}$ or $q \preceq \succ \tilde{q}$) iff there exists $p \in \text{Ps}(P, -T)$ or $\tilde{p} \in \text{Ps}(P, T)$ or $\tilde{\tilde{p}} \in \text{Ps}(P, T)$ such that $p \preceq q$ and $p \preceq \tilde{q}$ or $\tilde{p} \prec q$ and $\tilde{p} \prec \tilde{q}$ or $\tilde{\tilde{p}} \preceq q$ and $\tilde{\tilde{p}} \preceq \tilde{q}$, respectively.

3.17. Lemma. Let $q, \tilde{q} \in \text{Ss}(P, T)$. Then the following assertions hold:

- (i) $q \preceq \succeq \tilde{q}$ iff $q^- \prec \succ \tilde{q}^-$.
- (ii) $q \prec \succ \tilde{q}$ iff $q^+ \prec \succ \tilde{q}^+$.
- (iii) $q \preceq \succ \tilde{q}$ iff $q^+ \preceq \succ \tilde{q}^+$.

3.18. Lemma. Let $q, \tilde{q} \in \text{Ss}(P, T)$. Then the following three assertions are equivalent:

- (i) $q \preceq \succeq \tilde{q}$;
- (ii) $q^- \cap \tilde{q}^- \prec q^-, q^- \cap \tilde{q}^- \prec \tilde{q}^-$;
- (iii) $q^- \wedge \tilde{q}^- \prec q^-, q^- \wedge \tilde{q}^- \prec \tilde{q}^-$.

3.19. Lemma. Let $q, \tilde{q} \in \text{Ss}(P, T)$. Then the following three assertions are equivalent:

- (i) $q \prec \succ \tilde{q}$;
- (ii) $q^+ \cap \tilde{q}^+ \prec q^+, q^+ \cap \tilde{q}^+ \prec \tilde{q}^+$;
- (iii) $q^+ \wedge \tilde{q}^+ \prec q^+, q^+ \wedge \tilde{q}^+ \prec \tilde{q}^+$.

3.20. Lemma. Let $q, \tilde{q} \in \text{Ss}(P, T)$. Then the following three assertions are equivalent:

- (i) $q \preceq \succ \tilde{q}$;
- (ii) $q^+ \cap \tilde{q}^+ \preceq q^+, q^+ \cap \tilde{q}^+ \preceq \tilde{q}^+$;
- (iii) $q^+ \wedge \tilde{q}^+ \preceq q^+, q^+ \wedge \tilde{q}^+ \preceq \tilde{q}^+$.

3.21. Theorem. The positive local equivalence, the negative local equivalence and the local equivalence of symmetric pseudoprocesses in P over T are equivalence relations in the set $\text{Ss}(P, T)$.

Proof. See 3.17 and Theorem 5.15 in [5].

3.22. Theorem. Let $q, \tilde{q} \in \text{Ss}(P, T)$ have right local existence of solutions. Then $q \prec \succ \tilde{q}$ iff the following conditions are fulfilled:

- (i) $D_q = D_{\tilde{q}}, \mathcal{E}_q = \mathcal{E}_{\tilde{q}}$;

(ii) there exists a map

$$r^+ : L_q \rightarrow R$$

such that

$$r^+(s, u) > u \text{ for } u < \sup D_s, \quad r^+(s, u) = u \text{ for } u = \max D_s,$$

and

$$s|_{\langle u, r^+(s, u) \rangle} \in S_{\sim q}.$$

Proof follows from (1.6.3), 1.21, 1.14 (iii), 3.17 (ii) and Theorem 5.16 in [5].

3.23. Theorem. Let $q, \sim q \in \text{Ss}(P, T)$ have left local existence of solutions. Then $q \leq \cong \sim q$ iff the following conditions are fulfilled:

(i) $D_q = D_{\sim q}$, $\mathcal{S}_q = \mathcal{S}_{\sim q}$;

(ii) there exists a map

$$r^- : L_q \rightarrow R$$

such that

$$r^-(s, u) < u \text{ for } \inf D_s < u, \quad r^-(s, u) = u \text{ for } \min D_s = u$$

and

$$s|_{\langle r^-(s, u), u \rangle} \in S_{\sim q}.$$

Proof is similar to that of Theorem 5.16 in [5].

3.24. Theorem. Let $q, \sim q \in \text{Ss}(P, T)$ have bilateral local existence of solutions. Then $q \leq \cong \sim q$ iff the following conditions are fulfilled:

(i) $D_q = D_{\sim q}$, $\mathcal{E}_q = \mathcal{E}_{\sim q}$, $\mathcal{S}_q = \mathcal{S}_{\sim q}$;

(ii) there exist maps

$$r^+, r^- : L_q \rightarrow R$$

such that

$$r^+(s, u) > u \text{ for } \sup D_s > u, \quad r^+(s, u) = u \text{ for } \max D_s = u,$$

$$r^-(s, u) < u \text{ for } \inf D_s < u, \quad r^-(s, u) = u \text{ for } \min D_s = u$$

and

$$s|_{\langle r^-(s, u), r^+(s, u) \rangle} \in S_{\sim q}.$$

Proof follows from (1.6.3), 1.21, 1.14, 3.17 and Theorem 5.21 in [5].

3.25. Theorem. Let T be a closed subset of R and let $q, \sim q \in \text{S}(P, T)$ be solution complete processes. Then $q \leq \cong \sim q$ iff $q = \sim q$.

Proof follows from 3.17 (iii) and Theorem 5.21 in [5].